

Another introduction into radare2

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Overview

- Features
- Components
- Api examples
- Introduction into Esil
- Problems



Features

Radare2 is not just one tool or a conglomeration of several tools.

It provides libs for creating your own tools and there are bindings for many languages.

→ Even if you don't like programming in C, you can still use the R_API.



Features

- Disassembler/Assembler
- Analysis
- Debugger
- Base conversion
- Binary information
- Payload helpers
- ...



Components

- rax2 (base conversion)
- rasm2 (disassembler/assembler)
- rabin2 (binary information)
- radiff2 (binary diffing)
- ragg2/ragg2-cc (payload helper)
- rahash2 (hashing)
- radare2 (all comes together)



rax2

rax2 is a base converter:

hex to ternary:

```
$ rax2 Tx23  
1022t
```

bin to int:

```
$ rax2 101010d  
42
```



rasm2

rasm2 is an assembler and disassembler for a great number of archs:

Gameboy disassembler:

```
$ rasm2 -a gb -d ff  
rst 56
```

Show all supported archs:

```
$ rasm2 -L  
_d          8051      PD          8051 Intel CPU  
...  
...
```



rabin2

rabin2 provides information about binaries:

/bin/ls:

```
$ rabin2 -I /bin/ls
file/bin/ls
typeEXEC (Executable file)
pic false
canary true
nx true
crypto false
has_va true
```



rabin2

```
rootelf
class ELF64
lang c
arch x86
bits 64
machine AMD x86-64 architecture
os linux
subsys linux
endian little
strip true
static false
...
```



rabin2

resolve linked libs:

```
$ rabin2 -l /bin/ls          (lower case L)  
[Linked libraries]  
libcap.so.2  
libacl.so.1  
libc.so.6
```

3 libraries



radiff2

radiff2 can do binary diffing and provides creating patchfiles.

Examples? Not here, later.



rahash2

rahash2 can calculate hashes:

calculate sha1 hash of a string:

```
$ rahash2 -a sha1 -s radare2  
0x00000000-0x00000007 sha1: 0fc19fa3bbe77e97ece9f0e036444ee16b277a88
```

calculate md5 hash of a file:

```
$ rahash2 -a md5 /bin/r2  
/bin/r2: 0x00000000-0x000129d1 md5: 91fb7bf0b5bb8e81a2bc01dc17f4de16
```



ragg2/ragg2-cc

ragg2 helps creating payloads, and provides filters for IDS-circumvention:

create a tiny bin:

```
$ echo "int main() { write (1,\"hi\n\", 3); exit(0); }" > hi.c
$ ragg2-cc hi.c
$ ./hi.c.bin
hi
```



ragg2/ragg2-cc

use xor-encoder for circumvention:

```
$ ragg2 -e xor -c key=32 -B `ragg2-cc -x hi.c`  
6a2c596a205be8fffffffffc15e4883c60d301e48ffc6e2f9c92420202048492a209f212020  
68ad15d0dfdfdf9a2320202098212020202f25981c2020206010df2f25e3
```



radare2

radare2 is the hex-editor where everything comes together. It provides:

- different io-layers
- debugger
- basic code analysis
- esil-vm (wip)
- ... (no examples here :()



API-Examples

r_asm and r_anal api:

```
#include <r_asm.h>
#include <r_anal.h>
#include <r_types.h>
#include <stdio.h>

void main () {
    RAsm *a;
    RAsmOp *aop;
    RAnal *anal;
    RAnalOp *anop;
    ut8 hex = 0x00;
    a = r_asm_new ();
    anal = r_anal_new ();
    aop = R_NEW0(RAsmOp);
    anop = r_anal_op_new ();
    r_asm_use (a, "gb");
    r_anal_use (anal, "gb");
    r_anal_op (anal, anop, 0x0, &hex, 1);
    r_asm_set_pc (a, 0x0);
    r_asm_disassemble (a, aop, &hex, 1);
    printf ("0x%02x\t:t%s\t\\tcycle-length is %i\n", hex, aop->buf_asm, anop->cycles);
    free (aop);
    r_anal_op_free (anop);
    r_asm_free (a);
    r_anal_free (anal);
}
```



API-Examples

```
$ gcc $(pkg-config --libs --cflags r_asm r_anal) ex_asm_anal.c -o ex_asm_anal  
$ ./ex_asm_anal  
0x00:    nop -    cycle-length is 4
```



Introduction into esil

Esil is the language for the r2 internal vm.

- Syntax
- Esil basic operations
 - Custom Ops
- Internal vars
- Examples



Esil-Syntax

Esil is quite similar to forth:

- ',' as separator
- <src>,<dst>,op
- <src0>,<src1>,op
- <dst>,op



Esil basic operations

- mov → =
- cmp → ==
- neg → !
- mul → *
- add → +
- sub → -
- if → ?{
- else → }{
- read → []
- write → =[]
- ...



Esil custom ops

Sometimes esil basic ops are not sufficient for emulation. In that case it is possible to create a custom op in the analysis-plugin.

From libr/include/r_anal.h:

```
typedef int (*RAnalEsilOp)(RAnalEsil *esil);  
...  
R_API r_anal_esil_set_op (RAnalEsil *esil, const char *op, RAnalEsilOp code);
```



Esil internal vars

The vm provides useful information, such as carry, via internal vars. Those can be accessed with the prefix '%'. They are calculated on demand. Most of the esil basic ops store the old and the new value of the destination of the op in the esil-struct.

Internal vars:

- %z - checks if the new value is zero
- %r - cpu-regsize in bytes
- %p - parity of the new value
- %cx - checks if there was a carry from bit x
- %bx - checks if there was a borrow from bit x



Examples

Here are a few examples for the gameboy-z80:

jp 0x150	→	0x150,pc,=
cp 0x11	→	17,a,==,%z,Z,=%b4,H,=%b8,C,=,1,N,=
xor a	→	a,a,^=%z,Z,=,0,N,=,0,H,=,0,C,=
push de	→	2,sp,-=,de,sp,=[2]
pop hl	→	sp,[2],hl,=,2,sp,+=
ret Z	→	Z,{,sp,[2],pc,=,2,sp,+=,}
adc b	→	b,C,+ ,a,+=,%c3,H,=%c7,C,=,0,N,=
halt	→	HALT (custom op)



Problems

The screenshot shows a terminal window titled "Terminal — bash — 80x24". The command entered is:

```
[0x100002608 /bin/ls]> f tmp&&sr sp&&x 64&&dr=&&s-&&s tmp&&f-tmp&&pd
```

The output displays assembly code and memory dump information. The assembly code is:

```
offset    0 1 2 3 4 5 6 7 8 9 A B C D E F 0123456789ABCDEF
0x100304350 7e38 0000 b801 0000 00eb 05b8 ffff ffff ~8.....
0x100304360 c9c3 5548 89e5 5348 89f1 488b 5660 488b ..UH..SH..H.V`H.
0x100304370 4760 488b 5850 4839 5a50 7fld 7c22 488b G`H.XPH9ZP..|"H.
0x100304380 5858 4839 5a58 7f11 7c16 488d 7768 488d XXH9ZX..|.H.whH.
rax 0x0000000000000004    rbx 0x0000000000000004    rcx 0x00000001003042b0
rdx 0x0000000000000000    rdi 0x0000000100305220    rsi 0x0000000000000000
rbp 0x0000000100305260    rsp 0x0000000100304350    r8 0x000000010014bc74
r9 0x0000000000000000    r10 0x0000000000000000    r11 0x0000000000000000
r12 0x000000010014a81a   r13 0x0000000000000000    r14 0x0000000100304340
r15 0x0000000000000001   rip 0x0000000000000000    rflags =
0x100002608 0-             be380000b8  mov esi, 0xb8000038
0x10000260d 0              0100  add [rax], eax
0x10000260f 0              0000  add [rax], al
.=< 0x100002611 0           eb05  jmp 0x100002618 [1] | Looks like
| 0x100002613 0           b8fffffff mov eax, 0xffffffff | you need
`-> 0x100002618 0           c9  leave
0x100002619 0             c3  ret
; -----
0x10000261a 0              55  push rbp
0x10000261b 8+            4889e5  mov rbp, rsp
0x10000261e 8              488b5660  mov rdx, [rsi+0x60]
```

A vertical text annotation on the right side of the assembly code reads:

```
-----| Looks like
| you need
| some help
-----| oo _/
| | |
| | |
```

The terminal window has a dark background with light-colored text. The scroll bar on the right is blue.



Problems

Sometimes you run into a situation where static analysis fails, is not sufficient or it is just not possible to visualize things without confusing the users.

- call/jmp <reg>
- function-signatures
- gameboys halt-instruction



call/jmp <reg> (function-pointers)

Function-pointers are nice, if you are just a programmer and not interested in reverse-engineering. If you just analyse the opcode, you won't succeed. So you have to focus more on the context to reduce the number of possible destinations. Emulation can be great for this, even if it will take a lot of time to emulate all possible paths.

But how can you visualize all these context information, without confusing the user?

The esil-vm is a great example for this problem.



call/jmp <reg> (function-pointers)



call/jmp <reg> : esil-vm

The esil-vm right now consists of 26 basic ops and 67 extended ops. The op-parsing and execution happens in `r_anal_esil_parse()`, `runword()` and `iscommand()`.

So ... let's take a short look at the disassembly:

```
$ r2 /lib/libr_anal.so
```



call/jmp <reg> : esil-vm

```
[0x000df72a]> pd 25
      0x000df72a 0 0> e84dfeffff call sym.iscommand
      sym.iscommand()
      0x000df72f 0 0 85c0 test eax, eax
      0x000df731 0 0 744e je 0xdf781
      0x000df733 0 0 488b45f8 mov rax, [rbp-0x8]
      0x000df737 0 8+ 4885c0 test rax, rax
      0x000df73a 0 8 7445 je 0xdf781
      0x000df73c 0 8 488b45e8 mov rax, [rbp-0x18]
      0x000df740 0 16+ 488b8070010. mov rax, [rax+0x170]
      0x000df747 0 24+ 4885c0 test rax, rax
      0x000df74a 0 24 7426 je 0xdf772
      0x000df74c 0 24 488b45e8 mov rax, [rbp-0x18]
      0x000df750 0 32+ 488b8070010. mov rax, [rax+0x170]
      0x000df757 0 40+ 488b4de0 mov rcx, [rbp-0x20]
      0x000df75b 0 48+ 488b55e8 mov rdx, [rbp-0x18]
      0x000df75f 0 56+ 4889ce mov rsi, rcx
      0x000df762 0 64+ 4889d7 mov rdi, rdx
      0x000df765 0 72+ ff d0 call rax
      0x00000000() ; sym.imp._ITM_deregisterTMCloneTable
      0x000df767 0 72 85c0 test eax, eax
      0x000df769 0 72 7407 je 0xdf772
      0x000df76b 0 72 b801000000 mov eax, 0x1
      0x000df770 0 80+ eb6a jmp 0xdf7dc ; (sym.runword)
      LL ; JMP XREF from 0x000df74a (sym.runword)
      LL ; JMP XREF from 0x000df769 (sym.runword)
      LL -> 0x000df772 0 80 488b45f8 mov rax, [rbp-0x8]
      0x000df776 0 88+ 488b55e8 mov rdx, [rbp-0x18]
      0x000df77a 0 96+ 4889d7 mov rdi, rdx
      0x000df77d 0 104+ ff d0 call rax
      0x00000000() ; sym.imp._ITM_deregisterTMCloneTable
```



call/jmp <reg> : esil-vm

```
[0x000df57c]> pdf : CALL XREF from 0x000df72a (sym.runword)
= (fcn) sym.iscommand 170
    0x000df57c 0 0 55 push rbp
    0x000df57d 0 8+ 4889e5 mov rbp, rsp
    0x000df580 0 16+ 4881ecb0000 sub rsp, 0xb0
    0x000df587 0 20+ 4889bd68ffff mov [rbp-0x98], rdi
    0x000df58e 0 28+ 4889b560ffff mov [rbp-0xa0], rsi
    0x000df595 0 36+ 48899558ffff mov [rbp-0xa8], rdx
    0x000df59c 0 44+ 488b8560ffff mov rax, [rbp-0xa0]
    0x000df5a3 0 52+ 4889c7 mov rdi, rax
    0x000df5a6 0 60> e8f551faaff call sym.imp.sdb_hash
        sym.imp.sdb_hash(unk)
    0x000df5ab 0 60 89c1 mov ecx, eax
    0x000df5ad 0 68+ 488d8570ffff lea rax, [rbp-0x90]
    0x000df5b4 0 76+ ba10000000 mov edx, 0x10
    0x000df5b9 0 84+ 4889c6 mov rsi, rax
    0x000df5bc 0 92+ 4889cf mov rdi, rcx
    0x000df5bf 0 100> e80c47faaff call sym.imp.sdb_itoa
        sym.imp.sdb_itoa()
    0x000df5c4 0 100 488945f8 mov [rbp-0x8], rax
    0x000df5c8 0 108+ 488b8568ffff mov rax, [rbp-0x98]
    0x000df5cf 0 116+ 488b8058010 mov rax, [rax+0x158]
    0x000df5d6 0 124+ 488b55f8 mov rdx, [rbp-0x8]
    0x000df5da 0 132+ 4889d6 mov rsi, rdx
    0x000df5dd 0 140+ 4889c7 mov rdi, rax
    0x000df5e0 0 148> e8eb5cfaff call sym.imp.sdb_num_exists
        sym.imp.sdb_num_exists()
    0x000df5e5 0 148 85c0 test eax, eax
    0x000df5e7 0 148 7436 je 0xdf61f
    0x000df5e9 0 148 488b8568ffff mov rax, [rbp-0x98]
    0x000df5f0 0 156+ 488b8058010 mov rax, [rax+0x158]
    0x000df5f7 0 164+ 488b4df8 mov rcx, [rbp-0x8]
    0x000df5fb 0 172+ ba00000000 mov edx, 0x0
    0x000df600 0 180+ 4889ce mov rsi, rcx
    0x000df603 0 188+ 4889c7 mov rdi, rax
    0x000df606 0 196> e88542faaff call sym.imp.sdb_num_get
        sym.imp.sdb_num_get()
    0x000df60b 0 196 4889c2 mov rdx, rax
    0x000df60e 0 204+ 488b8558ffff mov rax, [rbp-0xa8]
    0x000df615 0 212+ 488910 mov [rax], rdx
    0x000df618 0 220+ b801000000 mov eax, 0x1
    0x000df61d 0 228+ eb05 jmp 0xdf624 ; (sym.iscommand)
: JMP XREF from 0x000df5e7 (sym.iscommand)
    0x000df61f 0 228 b800000000 mov eax, 0x0
: JMP XREF from 0x000df61d (sym.iscommand)
    0x000df624 0 236+ c9 leave
    0x000df625 0 228- c3 ret
```



call/jmp <reg> : esil-vm

The `is` command creates a hash of the op-string and checks if it exists in the sdb-instance. If so, it resolves a function-pointers belonging to the op-string. → `call rax`

- It's hard to resolve all possible values for `rax` at this point
- Even if you can resolve all possible values for `rax`, you cannot ensure if all of them will be used (custom ops)
- It's hard to visualize them, without spamming



Function-Signatures

Function-Signatures are not so good resolved by r2 at the moment.

If a function with 1 arg is called twice, r2 might show you a wrong signature



GB-Halt-Instruction

Halt on gameboy is nasty, because gameboy is nasty:

- checks if interrupts are enabled
- if so:
 - run the next **byte** twice (this can result in jr-trampolins)
 - endless-loop if the next byte is a halt-instruction
 - r2 cannot handle conditional repetition of only 1 byte
- if not:
 - skip the next **instruction**
 - easy for r2 to handle



Pointers

- <http://radare.org/>
- <http://radare.today/>
- <https://github.com/radare/radare2/>
- <irc://irc.freenode.net/radare>





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Reporting bugs



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memegenerator.net



EOF

- Questions?
- Ideas?

